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## The L(2, 1)-labeling on $\beta$ -product of Graphs

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**Abstract:** The L(2, 1)-labeling (or distance two labeling) of a graph G is an integer labeling of G in which two vertices at distance one from each other must have labels differing by at least 2 and those vertices at distance two must differ by at least 1. The L(2, 1)-labeling number of G is the smallest number k such that G has an L(2, 1)-labeling with maximum of f(v) is equal to k, where v belongs to vertex set of G. In this paper, upper bound for the L(2, 1)-labeling number for the  $\beta$ -product of two graphs has been obtained in terms of the maximum degrees of the graphs involved.

**Keywords:** Channel Assignment, L(2, 1)-labeling, L(2, 1)-labeling Number, Graph  $\beta$ -product

## 1. Introduction

The concept of L(2,1) -labeling in graph come into existence with the solution of frequency assignment problem. In fact, in this problem a frequency in the form of a nonnegative integer is to assign to each radio or TV transmitters located at various places such that communication does not interfere. Hale [6] was first person who formulated this problem as a graph vertex coloring problem. By Roberts [13], In order to avoid interference, any two "close" transmitters must receive different channels and any two "very close" transmitters must receive channels that are at least two channels apart. To translate the problem into the language of graph theory, the transmitters are represented by the vertices of a graph; two vertices are "very close" if they are adjacent and "close" if they are of distance two in the graph. Based on this problem, Griggs and Yeh [5] introduced L(2,1) -labeling on a simple graph. Let G be a graph with vertex set V(G). A function  $f:V(G)\to Z^+\cup\{0\}$ , where  $Z^+$  is a set of positive integers, is called L(2,1) -labeling or distance two labeling if  $|f(u) - f(v)| \ge 2$  when d(u, v) = 1and  $|f(u)-f(v)| \ge 1$  when d(u,v) = 2, where d(u,v) is distance between u and v in G. A k-L(2,1) -labeling is an L(2,1) -labeling such that no label is greater than k. The L(2,1) -labeling number of G, denoted by  $\lambda(G)$  or  $\lambda$ , is the smallest number k such that G has a k-L(2,1) -labeling. The L(2,1) -labeling has been extensively studied in recent past by many researchers [see, 1, 2, 4, 7, 8, 14-21]. The common trend in most of the research paper is either to determine the value of L(2,1) -labeling number or to suggest bounds for particular classes of graphs.

Griggs and Yeh [5] provided an upper bound  $\Delta^2 + 2\Delta$  for a general graph with the maximum degree  $\Delta$ . Later, Chang and Kuo [1] improved the upper bound to  $\Delta^2 + \Delta$ , while Kral and Skrekarski [10] reduced the upper bound to  $\Delta^2 + \Delta - 1$ . Furthermore, recently Gonccalves [4] proved the bound  $\Delta^2 + \Delta - 2$  which is the present best record. If G is a diameter 2 graph, then  $\lambda(G) \le \Delta^2$ . The upper bound is attainable by Moore graphs (diameter 2 graphs with order  $\Delta^2 + 1$ ). (Such graphs exist only if  $\Delta = 2, 3, 7$  and possibly 57 [5]). Thus Griggs and Yeh [5] conjectured that the best bound is  $\Delta^2$  for any graph G with the maximum degree  $\Delta \ge 2$ . (This is not true for  $\Delta = 1$ . For example,  $\Delta(K_2) = 1$  but  $\lambda(K_2) = 2$ ).

Graph products play an important role in connecting many useful networks. Klavzar and Spacepan [9] have shown that  $\Delta^2$ -conjecture holds for graphs that are direct or strong products of nontrivial graphs. After that Shao, et al. [15] have improved bounds on the L(2,1)-labeling number of direct and strong product of nontrivial graphs with refined approaches. Shao and Zhang [17] also consider the graph formed by the

Cartesian sum of graphs and prove that the  $\lambda$ -number of L(2,1)-labeling of this graph satisfies the  $\Delta^2$ -conjecture (with minor exceptions). Pradhan and Kumar [12] have obtained upper bound of the L(2,1)-labeling number for the  $\alpha$ -product of two graphs in terms of the maximum degree of the graphs involved and shown that the  $\lambda$ -number of L(2,1)

-labeling of this graph also satisfies the  $\Delta^2$  -conjecture (with minor exceptions)

In this paper, ithave considered the graph formed by the  $\beta$ -product of graphs [3] and obtained a general upper bound for L(2,1)-labeling number in term of maximum degree of the graphs. In the case of  $\beta$ -product of graphs, L(2,1)-labeling number of graph holds Griggs and Yeh's conjecture [5] with minor exceptions.

## 2. A labeling Algorithm

A subset X of V(G) is called an i-stable set (or i-independent set) if the distance between any two vertices in X is greater than i, i.e.  $\{d(u,v)>i, \ \forall \ u,v\in X\}$ . A 1-stable set is a usual independent set. A maximal 2-stable subset X of a set Y is a 2-stable subset of Y such that X is not a proper subset of any 2-stable subset of Y.

Chang and Kuo [1] proposed the following algorithm to obtain an L(2,1)-labeling and the maximum value of that labeling on a given graph.

Algorithm:

Input: A graph G = (V, E).

Output: The value k is the maximum label.

Idea: In each step i, find a maximal 2-stable set from the unlabeled vertices that are distance at least two away from those vertices labeled in the previous step. Then label all the vertices in that 2-stable with i in current stage. The label i starts from 0 and then increases by 1 in each step. The maximum label k is the final value of i.

Initialization: Set  $X_{-1} = \phi$ ; V = V(G); i = 0.

Iteration:

- 1) Determine  $Y_i$  and  $X_i$ .
  - a.  $Y_i = \{u \in V : u \text{ is unlabeled and } d(u, v) \ge 2 \ \forall v \in X_{i-1}\}$ .
  - b.  $X_i$  is a maximal 2-stable subset of  $Y_i$ .
  - c. If  $Y_i = \phi$  then set  $X_i = \phi$ .
- 2) Label the vertices of  $X_i$  (if there is any) with i.
- 3)  $V \leftarrow V X_i$ .
- 4) If  $V \neq \phi$ , then  $i \leftarrow i+1$ , go to step 1.
- 5) Record the current i as k (which is the maximum label). Stop.

Thus k is an upper bound on  $\lambda(G)$ .

Let u be a vertex with largest label k obtained by above Algorithm. We have the following sets on the basis of Algorithm just defined above.

 $I_1 = \{i : 0 \le i \le k-1 \text{ and } d(u,v) = 1 \text{ for some } v \in X_i\},\$ 

i.e.  $I_1$  is the set of labels of the neighbourhood of the vertex u.

 $I_2 = \{i : 0 \le i \le k-1 \text{ and } d(u,v) \le 2 \text{ for some } v \in X_i\},$ 

i.e.  $I_2$  is the set of labels of the vertices at distance at most 2 from the vertex  $\boldsymbol{u}$  .

$$I_3 = \{i : 0 \le i \le k-1 \text{ and } d(u,v) \ge 3 \text{ for all } v \in X_i\}$$
  
=  $\{0,1,...,k-1\} - I_2$ .

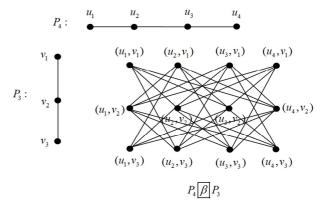
i.e.  $I_3$  consists of the labels not used by the vertices at distance at most 2 from the vertex u.

Then Chang and Kuo showed that  $\lambda(G) \le k \le |I_2| + |I_3| \le |I_2| + |I_1|$ .

In order to find k, it suffices to estimate  $B = |I_1| + |I_2|$  in terms of  $\Delta(G)$ . We will investigate the value B with respect to a particular graph ( $\beta$ -product of two graphs). The notations which have been introduced in this section will also be used in the following sections.

## 3. The $\beta$ -product of Graphs

The  $\beta$ -product  $G[\beta]H$  of two graphs G and H is the graph with vertex set  $V(G)\times V(H)$ , in which the vertex (u,v) is adjacent to the vertex (u',v') if and only if either  $u\neq u'$  and v is adjacent to v' in H or u is adjacent to u' in G and  $v\neq v'$  i.e. either  $u\neq u'$  and  $vv'\in E(H)$  or  $uu'\in E(G)$  and  $v\neq v'$ . For example, consider the Figure 1.



**Figure 1.**  $\beta$ -product of  $P_4$  and  $P_3$ .

## 4. The Maximum Degree (Largest Degree) of $G[\beta]H$

The maximum (largest) degree of  $G[\beta]H$  plays an important role in finding out the upper bound for the L(2,1)-labeling. To find out the maximum degree of  $G[\beta]H$ , we proceed as follows:

Let  $\Delta_1, \Delta_2$  be the maximum degree of G, H and  $\Delta'_1, \Delta'_2$  be the minimum degree of G, H respectively. Let  $\Delta$  be the

maximum degree of  $G[\overline{\beta}]H$  .

Case I: If  $n_1 \ge n_2$ , then

$$\Delta = \begin{cases} (n_1 - 1)\Delta_2 + (n_2 - 1)\Delta'_1 - \Delta'_1 \Delta_2 & \text{when } \Delta_1 \le \Delta_2 \\ (n_1 - 1)\Delta'_2 + (n_2 - 1)\Delta_1 - \Delta_1 \Delta'_2 & \text{when } \Delta_1 > \Delta_2 \end{cases}$$

Case II: If  $n_1 < n_2$ , then

$$\Delta = \begin{cases} (n_1 - 1)\Delta'_2 + (n_2 - 1)\Delta_1 - \Delta_1\Delta'_2 & \text{when } \Delta_1 \ge \Delta_2 \\ (n_1 - 1)\Delta_2 + (n_2 - 1)\Delta'_1 - \Delta'_1\Delta_2 & \text{when } \Delta_1 < \Delta_2 \end{cases}$$

From the above two cases, it can be written as

$$\Delta \leq \max \left[ (n_1 - 1)\Delta_2 + (n_2 - 1)\Delta_1' - \Delta_1' \Delta_2, (n_1 - 1)\Delta_2' + (n_2 - 1)\Delta_1 - \Delta_1\Delta_2' \right].$$

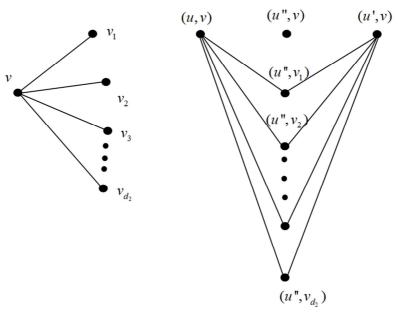


Figure 2. Internally-disjoint graph of paths.

# 5. Upper Bound for the L(2,1) – labeling Number in $G[\beta]H$

In this section, general upper bound for the L(2,1)-labeling number ( $\lambda$ -number) of  $\beta$ -product in term of maximum degree of the graphs has been established. In this regard, it state and prove the following theorem.

Theorem 5.1: Let  $\Delta$ ,  $\Delta_1$ ,  $\Delta_2$  be the maximum degree of  $G[\beta]H$ , G, H and n,  $n_1$ ,  $n_2$  be the number of vertices of  $G[\beta]H$ , G, H respectively. If  $\Delta_1$ ,  $\Delta_2 \ge 2$  , then  $\lambda(G[\beta]H) \le \Delta^2 + \Delta - (\Delta_1 - 1)(\Delta_2 - 1)(\Delta_1 + \Delta_2) - \Delta_1\Delta_2(n_1 + n_2 - \Delta_1 - \Delta_2 - 2) + 1$ . Proof: Let  $u_\beta = (u, v)$  be any vertex in the graph  $G[\beta]H$ . Denote  $d = \deg_{G[\beta]H}(u_\beta)$ ,  $d_1 = \deg_G(u)$ ,  $d_2 = \deg_H(v)$ ,  $\Delta_1 = \max \deg(G)$ ,  $\Delta'_1 = \min \deg(G)$ ,  $\Delta_2 = \max \deg(H)$ ,  $\Delta'_2 = \min \deg(H)$ ,  $|v(G)| = n_1$ , and  $|v(H)| = n_2$ . Hence it is easily shown that  $d = (n_1 - 1)d_2 + (n_2 - 1)d_1 - d_1d_2$  and  $\Delta = \Delta(G[\beta]H) \le \max[(n_1 - 1)\Delta_2 + (n_2 - 1)\Delta'_1 - \Delta'_1\Delta_2, (n_1 - 1)\Delta'_2 + (n_2 - 1)\Delta_1 - \Delta_1\Delta'_2]$ .

Consider the Figure 2. For any vertex u' in G at distance 2 from u, there must be a path u'u''u of length two between u' and u in G; but the degree of v in H is  $d_2$ , i.e. v has  $d_2$  adjacent vertices in H , by the definition of  $oldsymbol{eta}$  -product  $G[\beta]H$ , there must be  $d_2$  internally-disjoint paths(two paths are said to be internally-disjoint if they do not intersect each other) of length two between (u',v) and (u,v). Hence for any vertex in G at distance 2 from  $u_{\beta} = (u, v)$  which are coincided in  $G[\beta]H$ ; on the contrary whenever there is no such a vertex in G at distance 2 from u in G, there exist no such corresponding  $d_2$  vertices at distance 2 from  $u_\beta = (u, v)$ which are coincided in  $G[\beta]H$ . In the former case, since such  $d_2$  vertices at distance 2 from  $u_\beta = (u, v)$  are coincided in  $G[\beta]H$  and hence they should be counted once only and therefore we have to subtract  $(d_2-1)$  from the value  $d(\Delta-1)$  which is the best possible number of vertices at distance 2 from a vertex  $u_{\beta} = (u, v)$  in  $G[\beta]H$ . Let the number of vertices in G at distance 2 from u be t, then  $t \in [0, d_1(\Delta_1 - 1)]$ . Now, if we take  $t = d_1(\Delta_1 - 1)$  which is the best possible number of vertices at distance 2 from a vertex u in G, then to get the number of vertices at distance 2 from  $u_{\beta} = (u,v)$  in  $G[\underline{\beta}]H$ , we will have to subtract at least  $d_1(d_2-1)(\Delta_1-1)$  from the value  $d(\Delta-1)$ .

For H, it can proceed in similar way to get the number of vertices at distance 2 from  $u_{\beta}=(u,v)$  in  $G[\beta]H$  and in this case subtract  $d_2(d_1-1)(\Delta_2-1)$  from the value  $d(\Delta-1)$ . Hence, the number of vertices at distance 2 from  $u_{\beta}=(u,v)$  in  $G[\beta]H$  will decrease  $d_1(d_2-1)(\Delta_1-1)+d_2(d_1-1)(\Delta_2-1)$  from the value  $d(\Delta-1)$  altogether. By the above analysis, the number  $d(\Delta-1)-d_1(d_2-1)(\Delta_1-1)-d_2(d_1-1)(\Delta_2-1)$  is now the best possible number of vertices at distance 2 from  $u_{\beta}=(u,v)$  in  $G[\beta]H$ .

But some cases are remaining to be considered for finding out the best possible number of vertices at distance 2 from  $u_{\beta} = (u, v)$  in  $G[\beta]H$ .

Let  $\varepsilon$  be the number of edges of the subgraph F induced by the neighbours of  $u_{\beta}$ . The edges of the subgraph F induced by the neighbours of  $u_{\beta}$  can be divided into the following case.

Consider the Figure 3. For each neighbour  $(u',v_t)$  (where u' is adjacent to u in G and  $v_t$  is any vertex of H which is neither equal to v nor adjacent to v) and  $(u_t,v')$  (where  $u_t$  is any vertex of G which is neither equal to u nor adjacent to u and v' is adjacent to v in H) of  $u_{\beta} = (u,v)$  in  $G[\beta]H$ . Obviously vertex  $(u_t,v')$  is adjacent to  $(u',v_t)$  in  $G[\beta]H$ , hence there is an edge between them. There are at least  $(n_1-d_1-1)d_2$  neighbour vertices like  $(u_t,v')$  of  $u_{\beta} = (u,v)$  in  $G[\beta]H$  and at least  $d_1$  vertices like  $(u',v_t)$  in  $G[\beta]H$ . So at least  $(n_1-d_1-1)d_1d_2$  edges will exist between  $(u_t,v')$  and  $(u',v_t)$ . Hence the number of edges of the subgraph F induced by the neighbours of  $u_{\beta}$  is at least

 $(n_1-d_1-1)d_1d_2$ . By symmetric analysis, the neighbours of  $u_\beta$  should again add at least  $(n_2-d_2-1)d_1d_2-1$ .

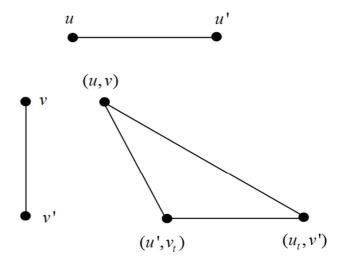


Figure 3. Combined graph of paths.

By the above analysis, we have

$$\varepsilon \ge (n_1 - d_1 - 1)d_1d_2 + (n_2 - d_2 - 1)d_1d_2 - 1$$
  
=  $(n_1 + n_2 - d_1 - d_2 - 2)d_1d_2 - 1$ .

Whenever there is an edge in F, the number of vertices with distance 2 from  $u_{\beta}$  will decrease by 2, hence the number of vertices with distance 2 from  $u_{\beta} = (u,v)$  in  $_{G}[\underline{\beta}]H$  will still need at least a decrease  $(n_1+n_2-d_1-d_2-2)d_1d_2-1$  from the value  $d(\Delta-1)-d_1(d_2-1)(\Delta_1-1)-d_2(d_1-1)(\Delta_2-1)$  (the number  $d(\Delta-1)-d_1(d_2-1)(\Delta_1-1)-d_2(d_1-1)(\Delta_2-1)$  is now the best possible for the number of vertices with distance 2 from  $u_{\beta} = (u,v)$  in  $G[\overline{\beta}]H$ ).

Hence for the vertex  $u_{\beta}$ , the number of vertices with distance 1 from  $u_{\beta}$  is no greater than  $\Delta$  and the number of vertices with distance 2 from  $u_{\beta}$  is no greater than

$$d(\Delta-1)-d_1(d_2-1)(\Delta_1-1)-d_2(d_1-1)(\Delta_2-1)-(n_1+n_2-d_1-d_2-2)d_1d_2+1.$$

Hence  $|I_1| \le d$ ,

$$|I_2| \le d + d(\Delta - 1) - d_1(d_2 - 1)(\Delta_1 - 1) - d_2(d_1 - 1)(\Delta_2 - 1) - (n_1 + n_2 - d_1 - d_2 - 2)d_1d_2 + 1$$

Then

$$B = |I_1| + |I_2|$$

$$= d + d + d(\Delta - 1) - d_1(d_2 - 1)(\Delta_1 - 1) - d_2(d_1 - 1)(\Delta_2 - 1) - (n_1 + n_2 - d_1 - d_2 - 2)d_1d_2 + 1$$

$$= d(\Delta + 1) - d_1(d_2 - 1)(\Delta_1 - 1) - d_2(d_1 - 1)(\Delta_2 - 1) - (n_1 + n_2 - d_1 - d_2 - 2)d_1d_2 + 1$$

$$= ((n_1 - 1)d_2 + (n_2 - 1)d_1 - d_1d_2)(\Delta + 1) - d_1(d_2 - 1)(\Delta_1 - 1) - d_2(d_1 - 1)(\Delta_2 - 1) - (n_1 + n_2 - d_1 - d_2 - 2)d_1d_2 + 1$$

Define

$$f(s,t) = ((n_1 - 1)t + (n_2 - 1)s - st)(\Delta + 1) - s(t - 1)(\Delta_1 - 1) - t(s - 1)(\Delta_2 - 1) - (n_1 + n_2 - s - t - 2)st + 1$$

Then f(s,t) has the absolute maximum at  $(\Delta_1, \Delta_2)$  on  $[0, \Delta_1] \times [0, \Delta_2]$ .

$$f(\Delta_1, \Delta_2) = ((n_1 - 1)\Delta_2 + (n_2 - 1)\Delta_1 - \Delta_1\Delta_2)(\Delta + 1) - \Delta_1(\Delta_2 - 1)(\Delta_1 - 1) - \Delta_2(\Delta_1 - 1)(\Delta_2 - 1)(\Delta_2 - 1)(\Delta_1 - 1) - \Delta_2(\Delta_1 - 1)(\Delta_2 - 1)(\Delta_2 - 1)(\Delta_1 - 1) - \Delta_2(\Delta_1 - 1)(\Delta_2 - 1)(\Delta_2 - 1)(\Delta_1 - 1)(\Delta_1 - 1)(\Delta_2 - 1)(\Delta_1 - 1)(\Delta_1 - 1)(\Delta_2 - 1)(\Delta_1 - 1$$

Since  $\Delta$  is the maximum degree of graph  $G[\beta]H$  and  $(n_1-1)\Delta_2 + (n_2-1)\Delta_1 - \Delta_1\Delta_2$  is the degree of any vertices in graph  $G[\beta]H$ . Therefore

$$(n_{1}-1)\Delta_{2} + (n_{2}-1)\Delta_{1} - \Delta_{1}\Delta_{2} \leq \Delta \leq \max \left[ (n_{1}-1)\Delta_{2} + (n_{2}-1)\Delta'_{1} - \Delta'_{1}\Delta_{2}, (n_{1}-1)\Delta'_{2} + (n_{2}-1)\Delta_{1} - \Delta_{1}\Delta'_{2} \right]$$

$$\therefore f(\Delta_{1}, \Delta_{2}) \leq \Delta(\Delta+1) - (\Delta_{2}-1)(\Delta_{1}-1)(\Delta_{1}+\Delta_{2}) - (n_{1}+n_{2}-\Delta_{1}-\Delta_{2}-2)\Delta_{1}\Delta_{2} + 1$$

$$= \Delta^{2} + \Delta - (\Delta_{2}-1)(\Delta_{1}-1)(\Delta_{1}+\Delta_{2}) - (n_{1}+n_{2}-\Delta_{1}-\Delta_{2}-2)\Delta_{1}\Delta_{2} + 1$$

Then,

$$\lambda(G[\beta]H) \le k \le B \le \Delta^2 + \Delta - (\Delta_2 - 1)(\Delta_1 - 1)(\Delta_1 + \Delta_2) - (n_1 + n_2 - \Delta_1 - \Delta_2 - 2)\Delta_1\Delta_2 + 1,$$

Where,  $\Delta \le \max [(n_1 - 1)\Delta_2 + (n_2 - 1)\Delta'_1 - \Delta'_1 \Delta_2, (n_1 - 1)\Delta'_2 + (n_2 - 1)\Delta_1 - \Delta_1 \Delta'_2].$ 

Corollary 5.2: Let  $\Delta$  be the maximum degree of  $G[\beta]H$ .

Then  $\lambda(G[\beta]H) \le \Delta^2 + \Delta - 3$  except for when one of  $\Delta(G)$  and  $\Delta(H)$  is 1.

Proof: If one of  $\Delta_1$  or  $\Delta_2$  is 1 then  $G[\beta]H$  is still a general graph, hence we can suppose that  $\Delta_1 \ge 2$  and  $\Delta_2 \ge 2$  (hence  $n_1 \ge 3$  and  $n_2 \ge 3$ ). Then

$$\Delta^{2} + \Delta - (\Delta_{2} - 1)(\Delta_{1} - 1)(\Delta_{1} + \Delta_{2}) - (n_{1} + n_{2} - \Delta_{1} - \Delta_{2} - 2)\Delta_{1}\Delta_{2} + 1$$

$$\leq \Delta^{2} + \Delta - (2 - 1)(2 - 1)(2 + 2) - (3 + 3 - 2 - 2 - 2).2.2 + 1$$

$$= \Delta^{2} + \Delta - 3$$

Therefore the result follows and holds Griggs and Yeh's conjecture [5] with minor exceptions.

## 6. Conclusion

In this paper, it has been discussed degree of vertex, maximum degree of vertex and number of vertices of maximum degree in  $\beta$ -product of two graphs. It has also find out the upper bound of L(2,1)-labeling number for  $\beta$ -product of two graphs in terms of the maximum degrees of the graphs involved. Upper bound of L(2,1)-labeling number for  $\beta$ -product of two graphs has been holds Griggs and Yeh's conjecture [5] with minor exceptions.

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